1. **Problem 1.** Suppose you are given a linear program

$$\max c^T x$$
$$Ax \le b$$
$$x \ge 0$$

And you know that x and y are primal and dual optimal. Then show that for any i either $(Ax)_i = b_i$ or $y_i = 0$.

2. **Problem 2.** Let \mathcal{P} be a polytope. Then recall that a point v is a vertex if we have that whenever $v = \frac{1}{2}(x+y)$ with $x, y \in \mathcal{P}$ then x = y = v. Prove that every polytope has a vertex.

(Hint: Consider the point $x \in \mathcal{P}$ that maximizes ||x||.)

- 3. Problem 3. For any polytope \mathcal{P} , prove the value of $c^T x$ is maximized on a vertex. (Note: there can be many points where the maximum point is achieved but at least 1 is a vertex solution.)
- 4. **Problem 4.** Consider a polytope given by $\mathcal{P} = \{x : Ax \leq b, x \geq 0\} \subset \mathbb{R}^d$. Prove that if x is a vertex of the polytope then we must have that at least d of the equations $Ax \leq b$ and $x \geq 0$ must hold with equality.
- 5. Problem 5. Prove that if the value of primal LP is unbounded then dual LP is unfeasible.
- 6. Problem 6. Formulate the following as an LP

$$\max |c^T x|$$
$$Ax \le b$$
$$x \ge 0$$

7. Problem 7. Formulate the following as an linear program.

$$\min ||Ax - b||_1$$

$$||x||_{\infty} \le 1$$

where $||x||_1 = \sum_i |x_i|$ and $||x||_{\infty} = \max_i |x_i|$.

(Source: http://www.seas.ucla.edu/ vandenbe/ee236a/homework/problems.pdf)

8. **Problem 8.** Consider the fractional vertex cover problem for a graph G = (V, E):

$$\min_{v \in V} x_v$$
$$\forall (u, v) \in E : x_u + x_v \ge 1$$
$$x_v \ge 0$$

Prove that there exists a half-integral optimal solution $x \in \{0, 1/2, 1\}^{|V|}$.

9. Problem 9. Extend the approximation algorithm given for vertex cover to one for a weighted vertex cover. That is for positive costs c_v find an approximate solution to:

$$\sum_{v \in V} c_v x_v$$
$$\forall (u, v) \in E : x_u + x_v \ge 1$$
$$x_v \in \{0, 1\}$$